BHCS15B: System Programming

Introduction to Compilers

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Course Web Page (www.mkbhandari.com/mkwiki)

Outline

Overview of Compilation

2 Phases of Compilation

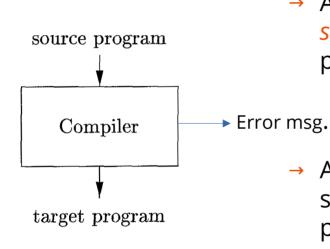
Introduction

 Programming languages are notations for describing computations to people and to machines.

- All the software running on all the computers are written in some programming language.
 - → The world depends on programming languages.
- A program must be translated into a form in which it can be executed by a computer.
 - → The software systems that do this translation are called compilers.

Language Processors

Compiler

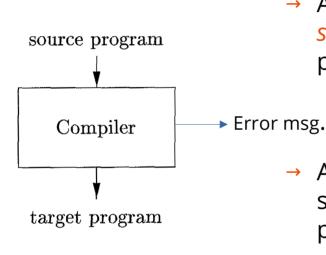


→ A program that can read a program in one language (*the source language*) and translates it into an equivalent program in another language (*the target language*).

→ An important role of a compiler is to *report errors* in the source program that it detects during the translation process.

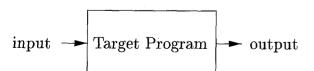
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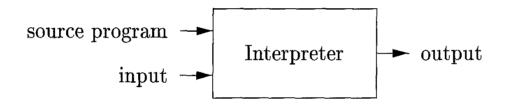


→ If the target program is an *executable machine-language program*, it can then be called by the user to process input and produce outputs.

Language Processors (2)

Interpreter

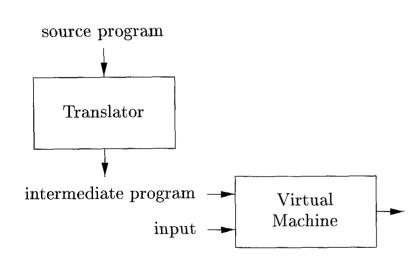
- → It is another common kind of language processor.
- → It appears to *directly execute the operations* specified in the source program on inputs supplied by the user



- → A compiler is usually *much faster* than an interpreter at mapping inputs to outputs.
- → An interpreter gives *better error diagnostics* than a compiler, because it executes the source program statement by statement.

Language Processors (3)

A Hybrid Compiler

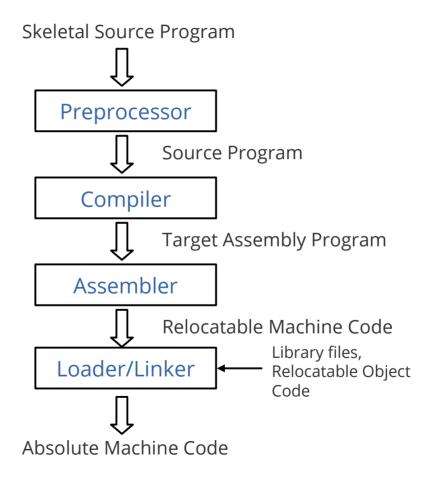


- → Java language processors combine compilation and interpretation
- → A Java source program are first be compiled into an intermediate form called bytecodes.
- → The *bytecodes* are then interpreted by a virtual machine.
- → A benefit of bytecodes is it is plateform independent (program compiled on one machine can be interpreted on another machine, perhaps across a network).

→ output

- → A just-in-time compilers in JVM helps to to achieve faster processing of inputs to outputs
- → Just-in-time compilers, translate the bytecodes into machine language immediately before they run the intermediate program to process the input.

Language Processors (4)



The structure of a Compiler

 So far, we have treated a compiler as a single box (unit) that maps a source program into semantically equivalent target program.

- It can be broadly divided into two parts:
 - Analysis (Front end of compiler)
 - Synthesis (Back end of compiler)

The structure of a Compiler (2)

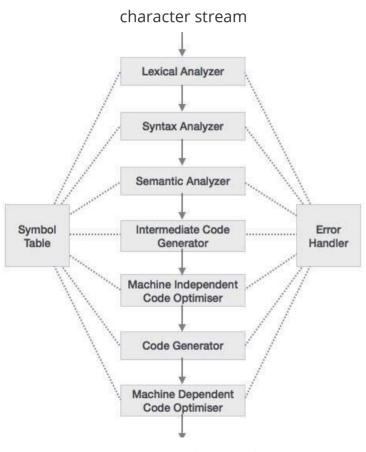
- Analysis part: breaks up the source program into constituent pieces and imposes a grammatical structure on them.
- Grammatical structure is then used to create an intermediate representation of the source program.
- If it detects that the source program is either syntactically ill formed or semantically unsound, then it must provide informative messages, so the user can take corrective action.
- It also collects information about the source program and stores it in the symbol table, which is passed along with the intermediate representation to the synthesis part.

The structure of a Compiler (3)

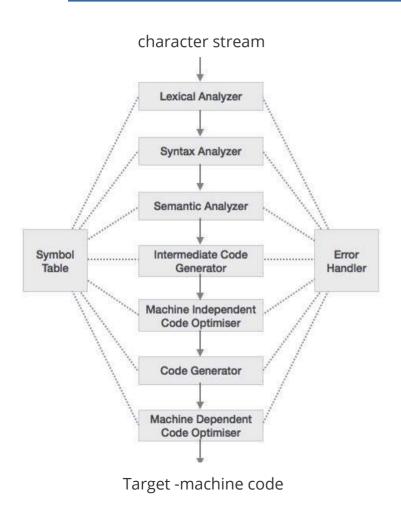
 Synthesis part: constructs the desired target program from the intermediate representation and the information in the symbol table.

• Compilation process: operates as a sequence of phases. Each phase takes input from its previous stage, has its own representation of source program, and feeds its output to the next phase of the compiler.

A typical decomposition of a compiler into phases is shown in next slide ->>



Target -machine code

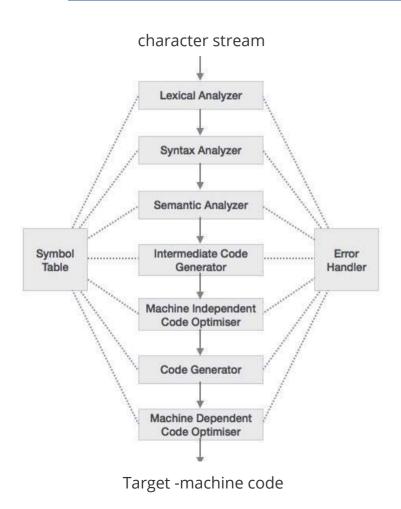


Lexical Analysis (Scanning):

- → The *lexical analyzer* reads the stream of characters (source program) and groups the characters into meaningful sequences called *lexemes*.
- → For each lexeme, the *lexical analyzer* produces as output a *token* of the form:

(token-name, attribute-value)

- → Tokens are then passed on to the next phase, i.e., syntax analysis.
- → Please see the example in Translation phase.jpg.

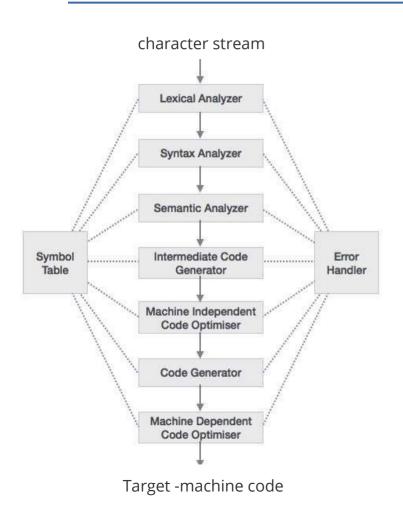


Syntax Analysis (Parsing):

→ The parser uses tokens produced by the *lexical* analyzer to create a syntax tree (an IR that depicts the grammatical structure of the token stream)

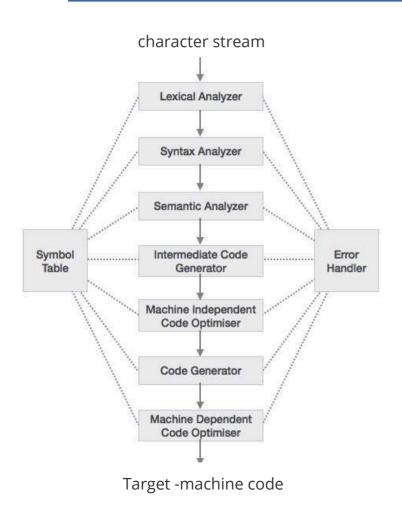
→ In syntax tree each interior node represents an operation and the children of the node represents the arguments of the operation.

Please see the example in Translation phase.jpg.



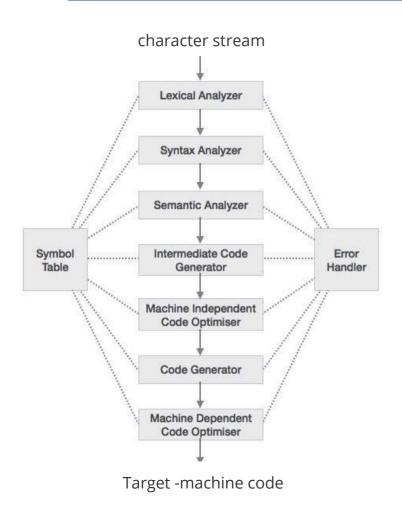
Semantic Analysis:

- → The semantic analyzer uses the *syntax tree* and the information in the symbol table to check the source program for semantic consistency with the language definition.
- → It also gathers type information and saves it either in the syntax tree or the symbol table, for subsequent use during intermediate-code generation.
- → *Type checking* is also done in this phase, where each operator is checked for matching operands.
- → Please see the example in Translation phase.jpg.



■ Intermediate Code Generation:

- → During compilation process, a compiler may construct one or more intermediate representations (IRs), which can have a variety of forms.
- → Syntax trees are a form of **IR**, commonly used during syntax and semantic analysis.
- → After syntax and semantic analysis, many compilers generate an explicit low-level or machine-like **IR** (program for an abstract machine).
- → This **IR** should be easy to produce and easy to translate into the target machine.
- → Please see the example in Translation phase.jpg.

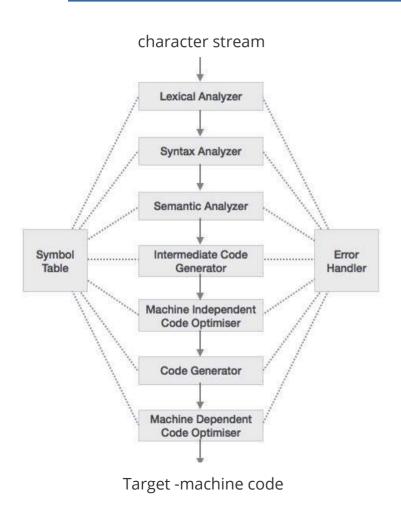


• *Code Optimization*:

- → The machine-independent code-optimization phase attempts to improve the intermediate code so that better target code will result.
 - faster
 - shorter

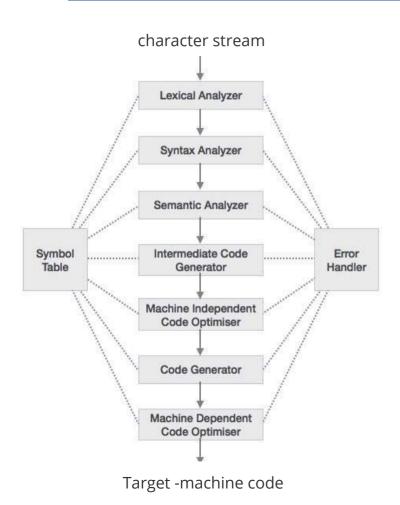
→ A simple intermediate code generation algorithm followed by code optimization is a reasonable way to generate good target code.

Please see the example in Translation phase.jpg.



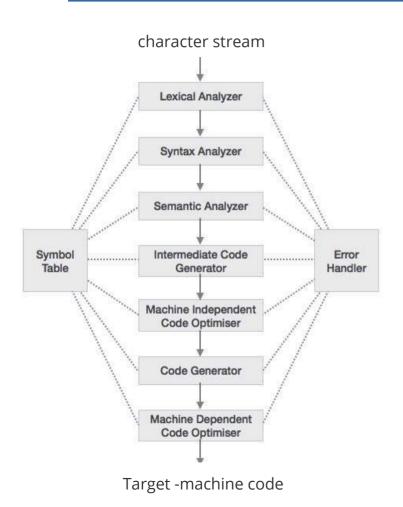
• Code Generation:

- → The code generator takes as input an **IR** of the source program and maps it into the target language.
- → If target language is machine code, then registers or memory locations are selected for each variables used by the program.
- → Then, the intermediate instructions are translated into sequence of machine instruction
- → Judicious assignment of the registers is the crucial aspect of code generation.
- → Please see the example in Translation phase.jpg.



Symbol Table Management

- → Records the variable names used in source program and collect information about their attributes:
 - Storage allocation, type, scope
- → In case of procedures names:
 - Number and type of arguments
 - Pass by value or pass by reference
 - Type returned
- → It should be designed such that, the store and retrieve operations are quick.
- Please see the example in Translation phase.jpg.



The grouping of Phases into Passes

- → In the implementation of a compiler, activities from several **phases** may be grouped together into a **pass** that reads an input file and writes an output file.
- → Front-end phases (lexical analysis, syntax analysis, semantic analysis, intermediate code generation) are grouped together into one pass:
- → Code optimization might be an optional pass.
- → Back-end pass consisting of code generation for a particular target machine.

References

Reference for this topic

- **Book:** Alfred V. Aho, Monica S. Lam, Ravi Sethi, J D Ullman, *Compilers: Principles, Techniques, and Tools*, 2nd Edition, Prentice Hall, 2006.
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